INFS 766 Internet Security Protocols

Lecture 1 Firewalls

Prof. Ravi Sandhu

INTERNET INSECURITY

- Internet insecurity spreads at Internet speed
 - > Morris worm of 1987
 - > Password sniffing attacks in 1994
 - > IP spoofing attacks in 1995
 - > Denial of service attacks in 1996
 - > Email borne viruses 1999
 - > Distributed denial of service attacks 2000
 - > Fast spreading worms and viruses 2003
 - > Spam 2004
 - > ... no end in sight
- * Internet insecurity grows at super-Internet speed
 - security incidents are growing faster than the Internet (which has roughly doubled every year since 1988)

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SECURITY OBJECTIVES

INTEGRITY disclosure AVAILABILITY access USAGE-CONTROL purpose 3

SECURITY TECHNIQUES

- Prevention
 - > access control
- Detection
 - > auditing/intrusion detection
 - > incident handling
- * Acceptance
 - > practicality

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THREATS, VULNERABILITIES ASSETS AND RISK

- * THREATS are possible attacks
- *** VULNERABILITIES are weaknesses**
- * ASSETS are information and resources that need protection
- RISK requires assessment of threats, vulnerabilities and assets

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RISK

***Outsider Attack**

insider attack

*Insider Attack

outsider attack

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PERSPECTIVE ON SECURITY

- * No silver bullets
- * A process NOT a turn-key product
- * Requires a conservative stance
- * Requires defense-in-depth
- * A secondary objective
- * Absolute security does not exist
- * Security in most systems can be improved

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PERSPECTIVE ON SECURITY

 absolute security is impossible does not mean absolute insecurity is acceptable

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INTRUSION SCENARIOS

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CLASSICAL INTRUSIONS SCENARIO 1

- * Insider attack
 - > The insider is already an authorized user
- Insider acquires privileged access
 - > exploiting bugs in privileged system programs
 - > exploiting poorly configured privileges
- Install backdoors/Trojan horses to facilitate subsequent acquisition of privileged access

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CLASSICAL INTRUSIONS SCENARIO 2

- Outsider attack
- Acquire access to an authorized account
- Perpetrate an insider attack

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NETWORK INTRUSIONS SCENARIO 3

- Outsider/Insider attack
- Spoof network protocols to effectively acquire access to an authorized account

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DENIAL OF SERVICE ATTACKS

- Flooding network ports with attack source masking
- *TCP/SYN flooding of internet service providers in 1996

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INFRASTRUCTURE ATTACKS

- * router attacks
 - > modify router configurations
- domain name server attacks
- * internet service attacks
 - > web sites
 - > ftp archives

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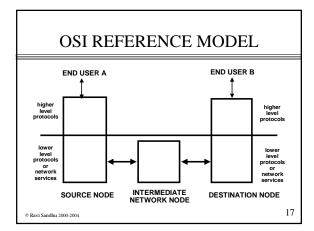
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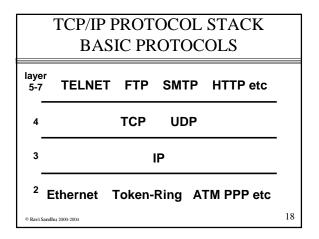
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INTERNET ARCHITECTURE AND PROTOCOLS

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OSI REFERENCE MODEL END USER A END USER B Application Laye Application Layer Presentation Layer Presentation Layer user functions Session Layer Session Layer Transport Layer Transport Layer Network Layer Network Layer Data Link Layer Data Link Layer Physical Layer Physical Layer PHYSICAL MEDIUM





TCP/IP PROTOCOL STACK BASIC PROTOCOLS

- **⋄ IP (Internet Protocol)**
 - > connectionless routing of packets
- *** UDP (User Datagram Protocol)**
 - > unreliable datagram protocol
- **❖ TCP (Transmission Control Protocol)**
 - > connection-oriented, reliable, transport protocol

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TCP/IP PROTOCOL STACK BASIC PROTOCOLS

- * TELNET: remote terminal
- * FTP (File Transfer Protocol)
- * TFTP (Trivial File Transfer Protocol)
- * SMTP (Simple Mail Transfer Protocol)
- * RPC (Remote Procedure Call)
- * HTTP (Hyper Text Transfer Protocol)
- * and others

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TCP/IP PROTOCOL STACK INFRASTRUCTURE PROTOCOLS layer TELNET FTP SMTP HTTP etc EGP DNS TCP UDP BGP ICMP IΡ ARP RARP Ethernet Token-Ring ATM PPP etc 21

TCP/IP PROTOCOL STACK INFRASTRUCTURE PROTOCOLS

- * ICMP: Internet Control Message Protocol
- * ARP: Address Resolution Protocol
- * RARP: Reverse Address Resolution Protocol
- * DNS: Domain Name Service
- * RIP: Routing Information Protocol
- * BGP: Border Gateway Protocol
- * EGP: External Gateway Protocol

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TCP/IP PROTOCOL STACK SECURITY PROTOCOLS layer TELNET FTP SMTP HTTP DNS SSL RIP EGP **TCP** UDP **BGP** IPSEC ICMP ΙP ARP RARP 2 Ethernet Token-Ring ATM 23

INTERNET STANDARDS PROCESS

- * IETF: Internet Engineering Task Force
 - > Application Area
 - > General Area
 - > Internet Area
 - > Operational Requirements Area
 - > Routing Area
 - > Security Area
 - > Transport Area
 - > User Services Area

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IETF SECURITY AREA ACTIVE WORKING GROUPS

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RFCs AND IETF DRAFTS

- * RFCs
 - > Standards
 - Proposed Standard
 - · Draft Standard
 - · Internet Standard
 - > Informational
 - > Experimental
 - > Historic
- IETF drafts
 - > work in progress
 - > expire after 6 months

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MUST, SHOULD, MAY

- * MUST
 - > mandatory, required of compliant implementations
- * SHOULD
 - > strongly recommended but not required
- * MAY
 - > possibility
 - > even if not stated a may is always allowed unless it violates MUST NOT

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TCP/IP VULNERABILITIES

BASIC TCP/IP **VULNERABILITIES**

- * many dangerous implementations of protocols
 - > sendmail
- * many dangerous protocols
 - > NFS, X11, RPC
 - > many of these are UDP based

BASIC TCP/IP **VULNERABILITIES**

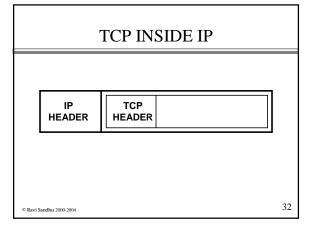
- * solution
 - > allow a restricted set of protocols between selected external and internal machines
 - > otherwise known as firewalls

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IP PACKET

- * header
- « data
 - > carries a layer 4 protocol
 - TCP, UDP
 - > or a layer 3 protocol
 - · ICMP, IPSEC, IP
 - > or a layer 2 protocol
 - · IPX, Ethernet, PPP

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IP HEADER FORMAT

- version: 4bit, currently v4
- * header length: 4 bit, length in 32 bit words
- * TOS (type of service): unused
- * total length: 16 bits, length in bytes
- identification, flags, fragment offset: total 16 bits used for packet fragmentation and reassembly
- * TTL (time to live): 8 bits, used as hop count
- Protocol: 8 bit, protocol being carried in IP packet, usually TCP, UDP but also ICMP, IPSEC, IP, IPX, PPP, Ethernet
- * header checksum: 16 bit checksum
- * source address: 32 bit IP address
- destination address: 32 bit IP address

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IP HEADER FORMAT

* options

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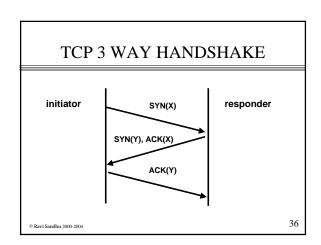
- > source routing
 - enables route of a packet and its response to be explicitly controlled
- > route recording
- > timestamping
- > security labels

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TCP HEADER FORMAT

- * source port number
 - > source IP address + source port number is a socket: uniquely identifies sender
- * destination port number
 - > destination IP address + destination port number is a socket : uniquely identifies receiver
- * SYN and ACK flags
- * sequence number
- * acknowledgement number

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TCP SYN FLOODING ATTACK

- * TCP 3 wav handshake
 - > send SYN packet with random IP source address
 - > return SYN-ACK packet is lost
 - this half-open connection stays for a fairly long time out period
- * Denial of service attack
- * Basis for IP spoofing attack

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IP SPOOFING

- Send SYN packet with spoofed source IP address
- SYN-flood real source so it drops SYN-ACK packet
- guess sequence number and send ACK packet to target
 - > target will continue to accept packets and response packets will be dropped

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TCP SESSION HIJACKING

- Send RST packet with spoofed source IP address and appropriate sequence number to one end
- * SYN-flood that end
- send ACK packets to target at other end

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SMURF ATTACK

- Send ICMP ping packet with spoofed IP source address to a LAN which will broadcast to all hosts on the LAN
- Each host will send a reply packet to the spoofed IP address leading to denial of service

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ULTIMATE VULNERABILITY

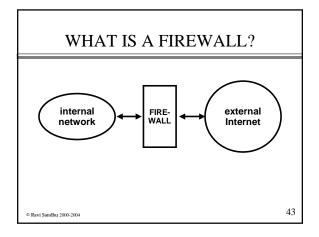
- IP packet carries no authentication of source address
- * IP spoofing is possible
 - > IP spoofing is a real threat on the Internet
 - IP spoofing occurs on other packet-switched networks also, such as Novell's IPX
- * Firewalls do not solve this problem
- * Requires cryptographic solutions

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FIREWALLS

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WHAT IS A FIREWALL?

- all traffic between external and internal networks must go through the firewall
 - > easier said than done
- firewall has opportunity to ensure that only suitable traffic goes back and forth
 - > easier said than done

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ULTIMATE FIREWALL internal network Air external Internet o Ravi Sandhu 2000-2004 45

BENEFITS

- secure and carefully administer firewall machines to allow controlled interaction with external Internet
- internal machines can be administered with varying degrees of care
- * does work

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BASIC LIMITATIONS

- * connections which bypass firewall
- services through the firewall introduce vulnerabilities
- insiders can exercise internal vulnerabilities
- * performance may suffer
- * single point of failure

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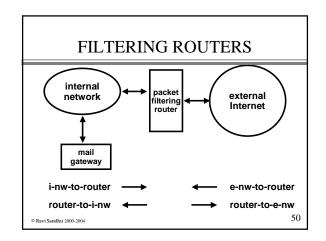
TYPES OF FIREWALLS

- Packet filtering firewalls
 - > IP layer
- * Application gateway firewalls
 - > Application layer
- Circuit relay firewalls
 - > TCP layer
- * Combinations of these

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PACKET FILTERING FIREWALLS

- * drop packets based on filtering rules
- * static (stateless) filtering
 - > no context is kept
- * dynamic (statefull) filtering
 - > keeps context

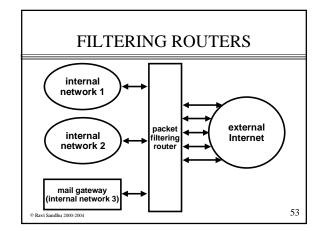
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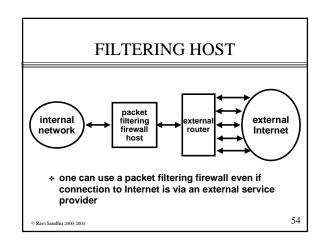
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PACKET FILTERING FIREWALLS

- Should never allow packet with source address of internal machine to enter from external internet
- Cannot trust source address to allow selective access from outside

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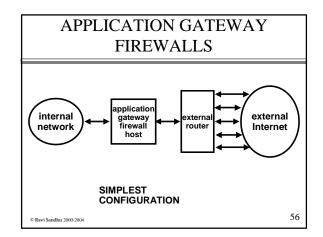
PACKET FILTERING FIREWALLS

- packet filtering is effective for coarse-grained controls
- not so effective for fine-grained control
 - > can do: allow incoming telnet from a particular host
 - > cannot do: allow incoming telnet from a particular user

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APPLICATION PROXIES

- have to be implemented for each service
- may not be safe (depending on service)

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CLIENT-SIDE PROXIES Internal-Client External-Server

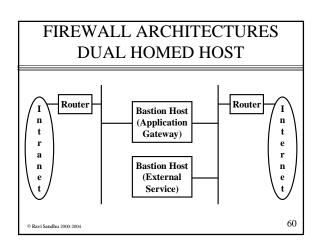
- allow outgoing http for web access to external machines from internal
- * requires some client configuration

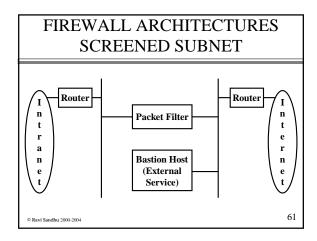
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SERVER-SIDE PROXIES External-Client Internal-Server

- allow incoming telnet for access to selected internal machines from selected external users
- requires some cryptographic protection to thwart sniffing and IP spoofing
- * becoming increasingly important for
- > electronic commerce
- > VPN
- > remote access security

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INTRUSION DETECTION

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RELATED TECHNOLOGIES

- * Intrusion detection
- Vulnerability assessment
- * Incident response
- * Honey pots
- Sniffer probes

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INTRUSION DETETCION TECHNIQUES

- * Policy detection (or knowledge-based)
 - > default permit
 - attack-signature based detection
 - also called misuse detection
 - > default deny
 - · specification-based detection
- * Anomaly detection (or behavior-based)
 - · requires user profiling
 - requires some learning capability in the system
- * Combinations of these

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INTRUSION DETECTION DATA SOURCE

- * network-based intrusion detection
 - > multiple sensor points
- * host-based intrusion detection
 - > multi-host based
- * application-based intrusion detection
- * combinations of these

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ATTACKER

- Outsider
 - > easier
- * insider
 - > harder

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INTRUSION DETECTION ISSUES

- * effectiveness
- * efficiency
- * security
- inter-operability
- * ease of use
- * transparency

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INTRUSION DETECTION CHALLENGES

- Performance and scalability

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BASE RATE FALLACY

- Test for a disease is 99% accurate
 - > 100 disease-free people tested, 99 test negative
 - > 100 diseased people tested, 99 test positive
- * Prevalence of disease is 1 in 10,000
- Alice tests positive
- * What is probability Alice has the disease?

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BASE RATE FALLACY

- ❖ Test for a disease is 99% accurate
 - > 100 disease-free people tested, 99 test negative
 - > 100 diseased people tested, 99 test positive
- ❖ Prevalence of disease is 1 in 10,000
- Alice tests positive
- What is probability Alice has the disease?
 1 in 100
- * False alarm rate: 99 in 100 !!!!!

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BASE RATE FALLACY BAYE'S THEOREM

⇒ population: 1,000,000

* diseased: 100

disease free: 999,900false positive: 9,999true positive: 99

* Alice's chance of disease: 99/(9,999+99) = 1/100

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BASE RATE FALLACY 99.99% ACCURACY

⋄ population: 1,000,000

* diseased: 100

disease free: 999,900false positive: 99.99true positive: 99.99

* Alice's chance of disease: 99.99/(99.99+99.99) = 1/2

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NETWORK-BASED INTRUSION DETECTION SIGNATURES

- * port signatures
- * header signatures
- string signatures

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NETWORK-BASED INTRUSION DETECTION ADVANTAGES

- Complements firewalls
- * broad visibility into network activity
- * no impact on network performance
- transparent installation

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NETWORK-BASED INTRUSION DETECTION DISADVANTAGES

- * False positives
- * miss new unknown attacks
- * scalability with high-speed networks
- * passive stance
- * emergence of switched Ethernet

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HOST-BASED INTRUSION DETECTION

- * host wrappers or personal firewalls
 - look at all network packets, connection attempts, or login attempts to the monitored machine
 - example, tcp-wrapper
- * host-based agents
 - monitor accesses and changes to critical system files and changes in user privilege
 - example, tripwire

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INTRUSION DETECTION STANDARDS

- * None exist
- ongoing efforts
 - > CIDF: common intrusion detection framework for sharing information
 - > IETF Intrusion Detection Working Group just started

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INTRUSION DETECTION

- Needs to integrate with other security technologies such as cryptography and access control
- one component of defense-in-depth layered security strategy
- incident-response and recovery are important considerations

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