



Label-Based Access Control: An ABAC Model with Enumerated Authorization Policy

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Summary

- Background & motivation
- Enumerated authorization policy ABAC model
- Relationship with existing models
- Expressive power of LaBAC
- Conclusion





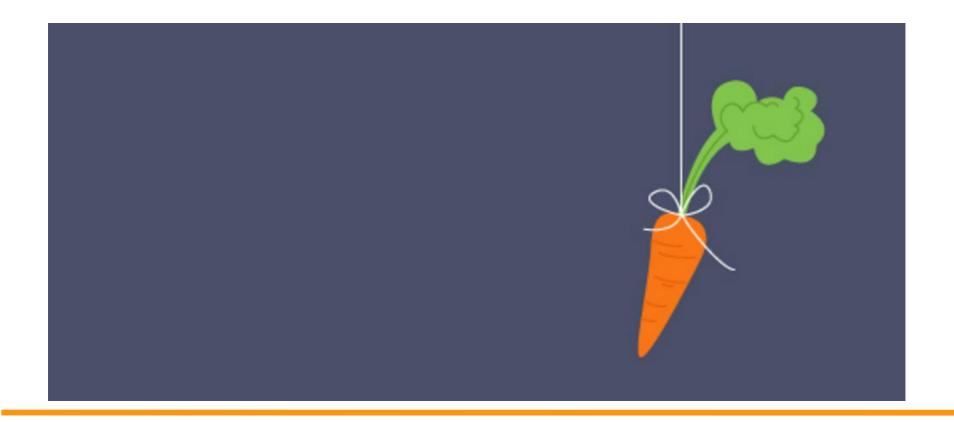
We present an enumerated authorization policy ABAC model and understand its relationship with traditional access control models.





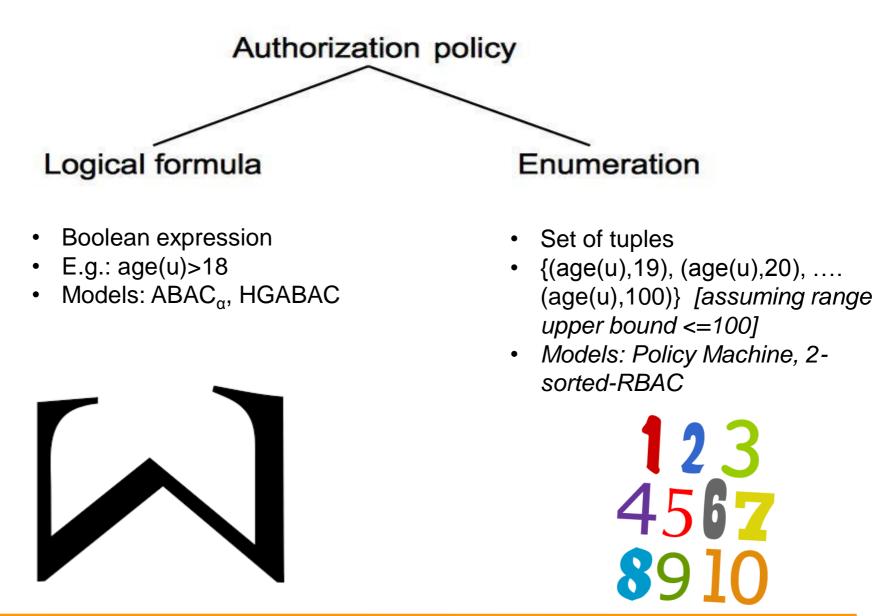


Background and Motivation













Many ways to set up a policy - **Auth**_{read} (**Auth**_{read} allows manager to read TS objects from home or office).

(i) $mng \in role(u) \land (office \in location(u) \lor home \in location(u)) \land TS \in sensitivity(o)$ (ii) $((mng \in role(u) \land office \in location(u)) \lor (mng \in role(u) \land home \in location(u))) \land TS \in sensitivity(o)$ (iii) $((mng \in role(u) \land office \in e location(u) \land TS \in sensitivity(o)) \lor ((mng \in role(u) \land home \in location(u) \land TS \in sensitivity(o)))$







Update **Auth**_{read} so that manager can no longer read TS objects from home

(i) $mng \in role(u) \land (office \in location(u) \lor home \in location(u)) \land TS \in sensitivity(o)$ (ii) $((mng \in role(u) \land office \in location(u)) \lor (mng \in role(u) \land home \in location(u)))$ $\land TS \in sensitivity(o)$ (iii) $((mng \in role(u) \land office \in e location(u) \land TS \in sensitivity(o)) \lor$ $((mng \in role(u) \land home \in location(u) \land TS \in sensitivity(o)))$







$\Box Auth_{read \equiv} \{(mng, home, TS), (mng, office, TS)\}$

□ Auth`_{read =} { (mng, home, TS), (mng,office,TS)}



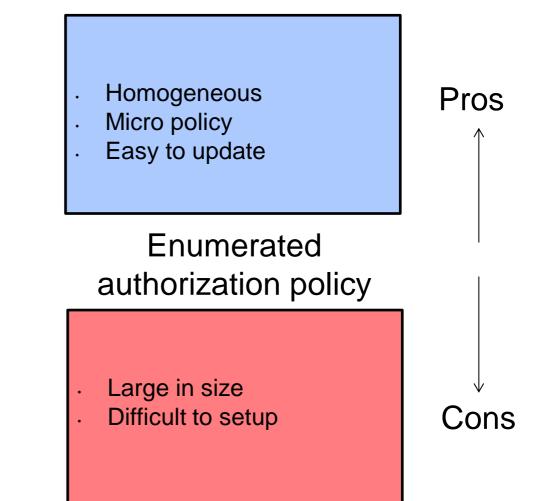




- · Rich & flexible
- . Easy to setup
- · Concise

Logical formula authorization policy

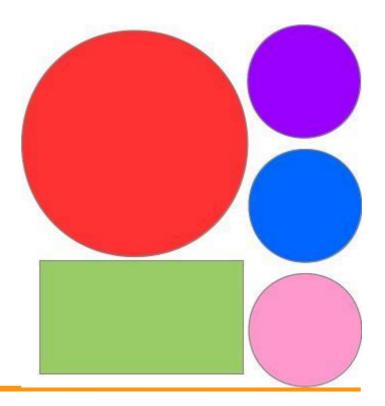
- Difficult to update
- Monolithic
- · Heterogeneous







LaBAC: Label-Based Access Control





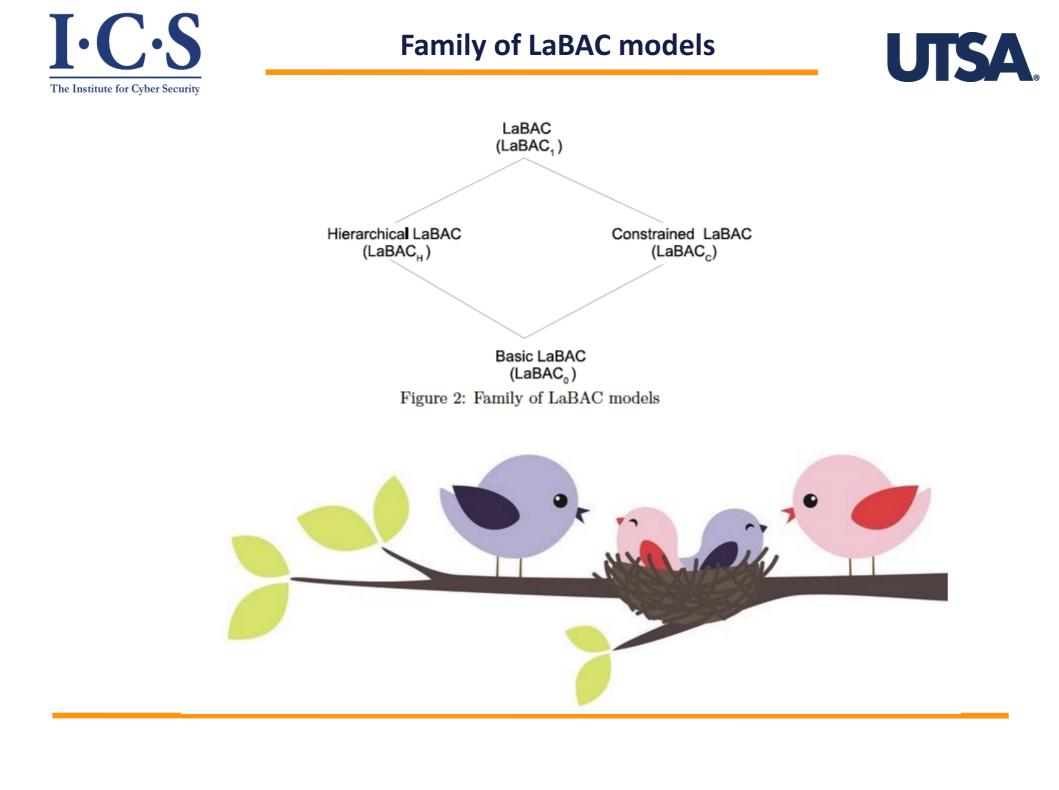


Label vs Attribute

Labels are attributes with tighter semantics

Salient features of LaBAC

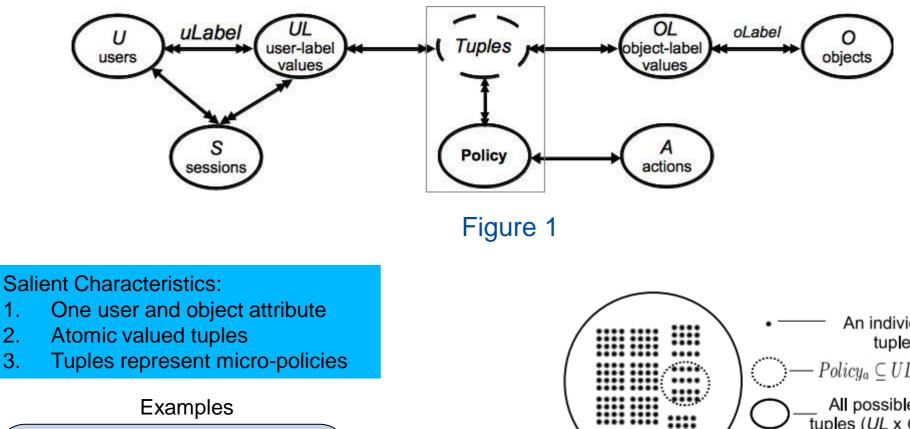
- Finite domain ABAC
- Simple enumerated ABAC model





LaBAC: Core model





UL={manager,employee}

 $OL=\{TS,S\}$

1.

2.

3.

Tuple1= (manager,TS)

Policy_{read} = {tuple1, tuple2...}

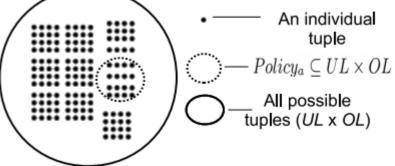
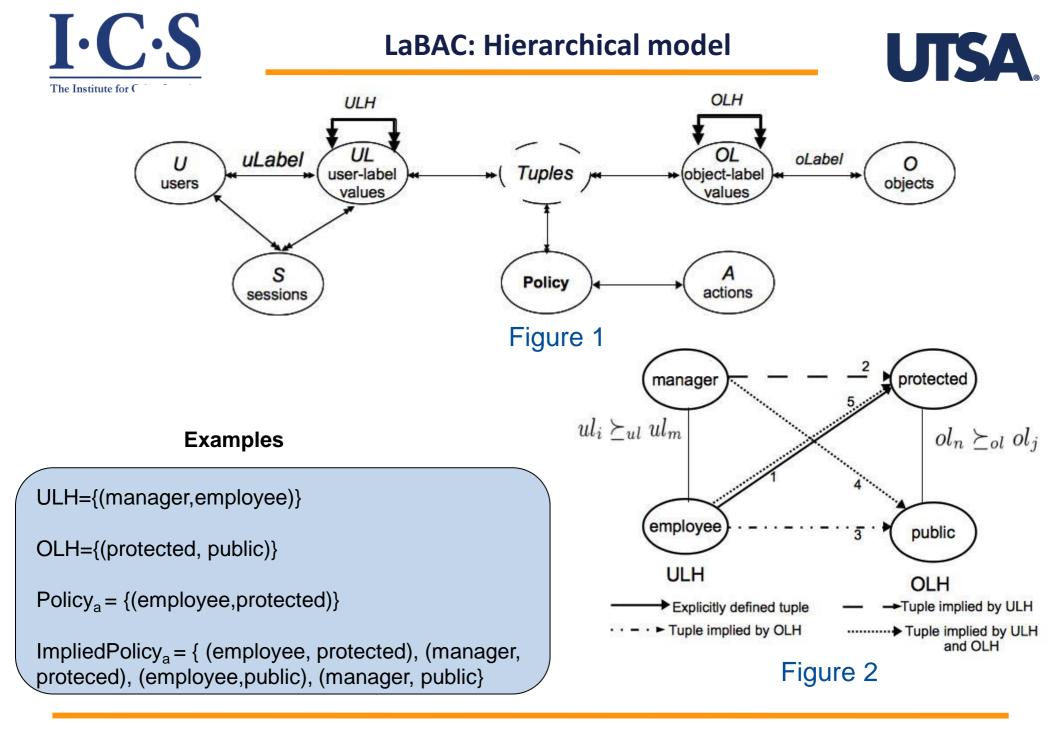
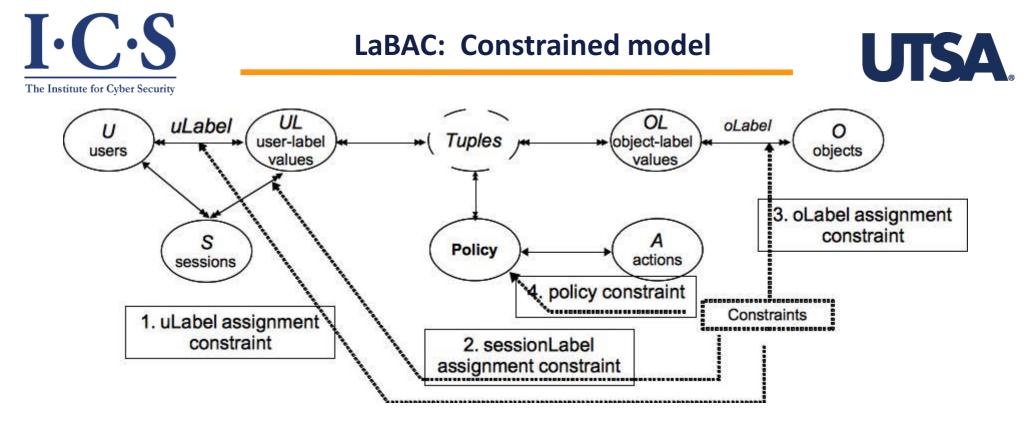


Figure 2





Examples

uLabel assgn. cons: a user cannot be both manager & director.

Session assgn. cons: at most one value can be activated in a session.

oLabel assgn. cons: A object cannot be both private & public

Policy cons: (employee, TS) can never be used.



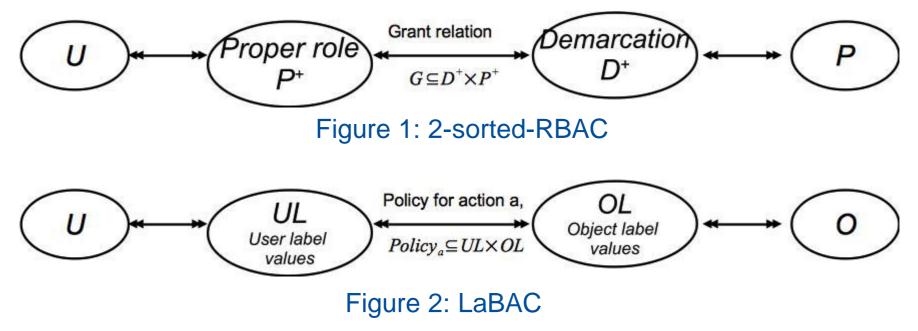


Relationship of LaBAC with other enumerated policy models









2-sorted-RBAC vs LaBAC:

- 1. Use of attributes
- 2. Separation of object and action from permission





□ Policy Machine _{mini}

- Only ASSIGN and ASSOCIATION relation
- Default policy class

□ Configuration of LaBAC in Policy Machine mini





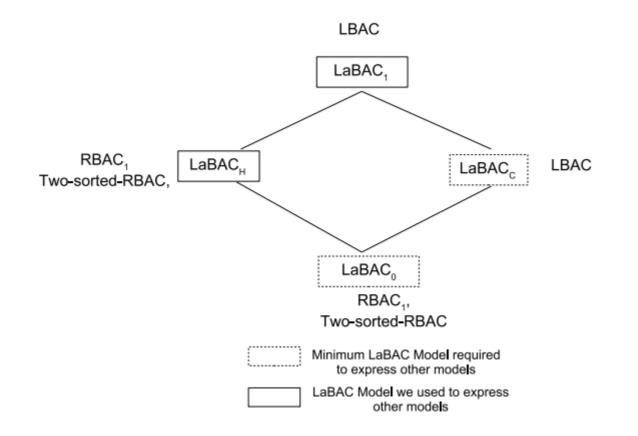


Flexibility in expressing traditional models





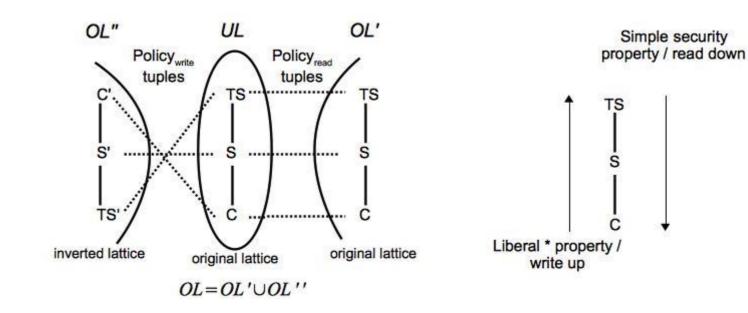






LBAC in LaBAC





LBAC assumptions:

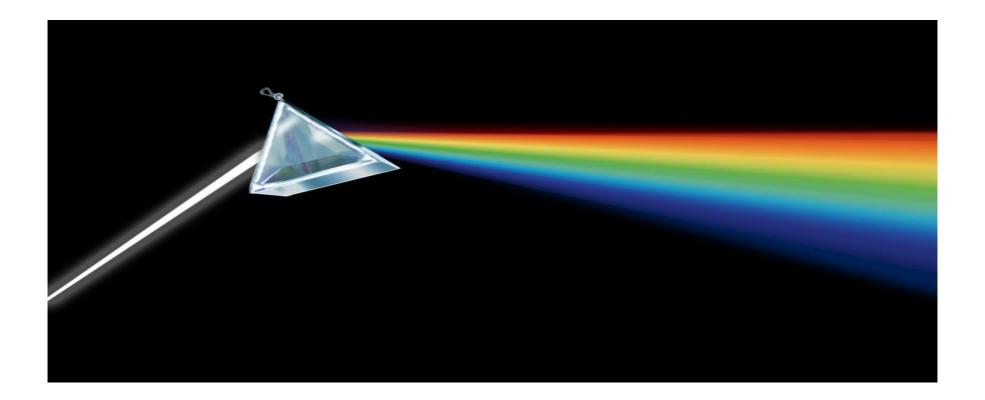
- 1. Tranquility
- 2. Object operation: creation only

- 1	
	UL = SC and $ OL = 2 * SC $
	Policy - 2 (Policy and Policy)
	$ FOUCY = 2 (FOUCY_{read} and FOUCY_{write}) $
	$ Policy = 2 \ (Policy_{read} \ \text{and} \ Policy_{write})$





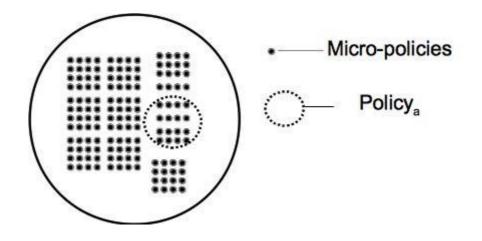
Micro-policy in LaBAC





Micro-policy in LaBAC





□ micro-policy as the smallest unit of administration

□ Example of a micro-policy: (manager, TS)





- □ Any other form of representation for authorization policy?
- □ How expressive power of enumerated authorization policy is compared with that of logical-formula auth. policy?
- What would be the cost of storing large number of enumerated tuples?













