

Institute for Cyber Security



Object-to-Object Relationship Based Access Control: Model and Multi-Cloud Demonstration

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Outline

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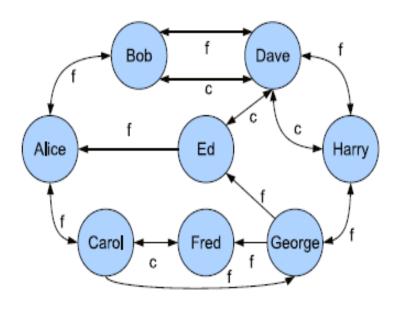






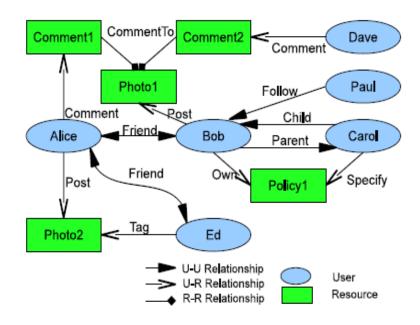
Relationships in OSN





User to user relationships in a sample social graph [UURAC,

Cheng et al. 2012

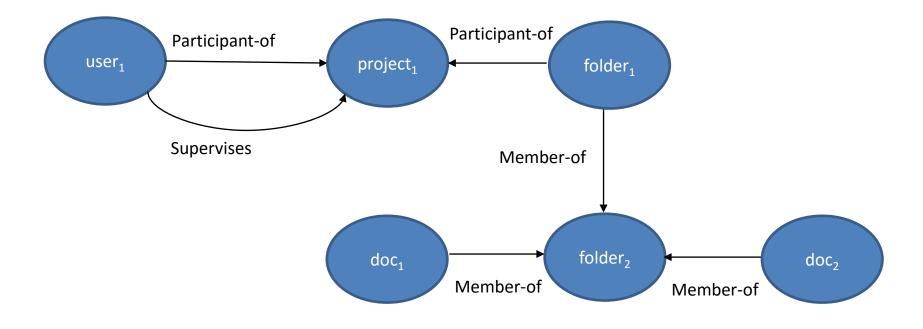


User to user, user to resource and resource to resource relationships in a sample social graph [URRAC, Cheng et al. 2012]



ReBAC in General Computing System beyond OSN





A sample Relationship Graph for Organizational Environment [RPPM, Crampton et al.,2014]



Object Relationships in Existing ReBAC



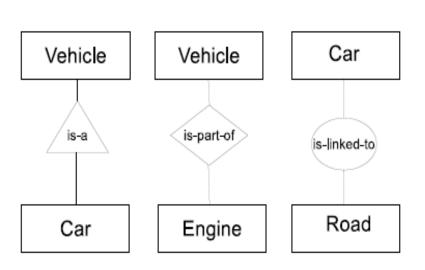
- ➤ Most of the ReBAC for OSN considers only user to user relationship
- ➤ OSN has very specific types of resources photos, notes, comments. Which only makes sense along with users.
- Even though some ReBAC models consider general computing system they still need users/subjects existence in relationship graph.

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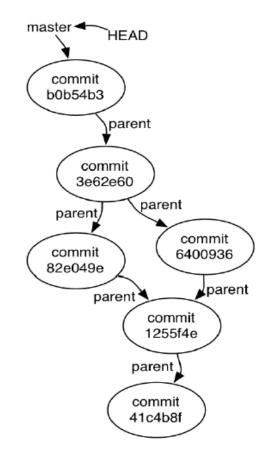


Existence of Object RelationshipIndependent of User





Object Relationship in Object –Oriented System (Inheritance, Composition and Association)



History of a Git Project (Version Control System) is a DAG



Limitations of Existing ReBAC Models



- Cannot configure relationship between objects independent of user.
- Cannot express authorization policy solely considering object relationship.

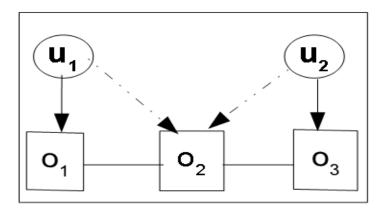
Considering these limitations we are proposing an object-to-object relationship based access control model.

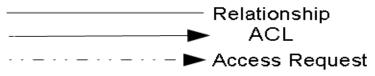


How the model would look like?



An Object to Object Relationship Based Access Control

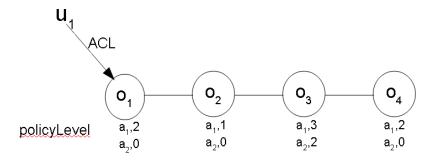




$$ACL(o_1) = \{u_1\}$$

 $ACL(o_2) = \{\}$
 $ACL(o_3) = \{u_2\}$

Policy Level Example

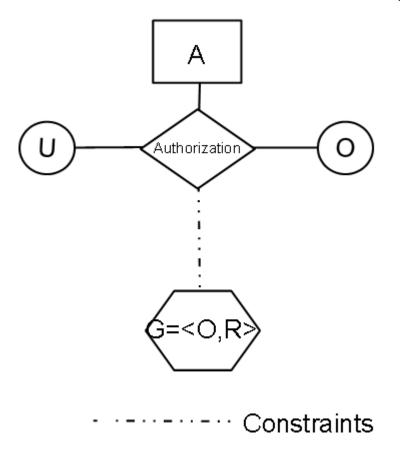


policyLevel $(a_1, o_1) = 2$ policyLevel $(a_2, o_1) = 0$ policyLevel $(a_1, o_2) = 1$ policyLevel $(a_2, o_2) = 0$ policyLevel $(a_1, o_3) = 3$ policyLevel $(a_2, o_3) = 2$ policyLevel $(a_1, o_4) = 2$ policyLevel $(a_2, o_4) = 0$



OOReBAC: Model Components and Definition





- U is a set of users
- O is a set of objects
- $R \subseteq \{z \mid z \subset O \land |z| = 2\}$
- G=(O,R) is an undirected relationship graph with vertices O and edges R
- · A is a set of actions
- Pⁱ(o₁) = { o₂ | there exists a simple path of length p in graph G from o₁ to o₂}
- policyLevel: × A → N
- ACL: O → 2^U which returns the Access control List of a particular object.
- There is a single policy configuration point. Authorization Policy, for each action a ∈ A, Authz_a(u:U,o:O) is a boolean function which returns true or false and u and o are formal parameters.
- Authorization Policy Language:
 Bach action "a" has a single authorization policy Authz_α(u:U,o:O) specified using the following language.
 φ := u ∈ PATH_i
 PATH_i := ACL(P⁰(o)) ∪ ... ∪ ACL(Pⁱ(o)) where i = min(| O | -1, policyLevel(a,o))

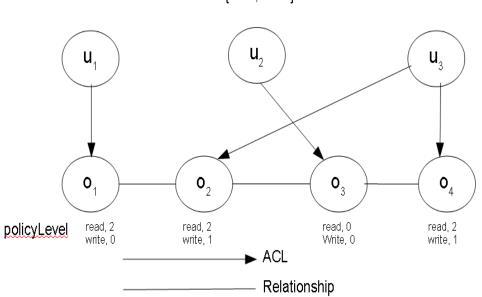
where for any set X, $ACL(X) = \bigcup_{x \in X} ACL(x)$



OOReBAC: An Example



A = {read, write}



Configuration:

- A = {read, write}
- $\begin{array}{l} \bullet \;\; Authz_{\mathit{read}}(u : U, o : O) \equiv u \in P^{\mathit{policyLevel(read,o)}} \\ \bullet \;\; Authz_{\mathit{write}}(u : U, o : O) \equiv u \in P^{\mathit{policyLevel(write,o)}} \end{array}$

Sequence of operations and its outcome:

- $U = \{u_1, u_2, u_3\}$
- $0 = \{ o_1, o_2, o_3, o_4 \}$
- $R = \{\{o_1, o_2\}, \{o_2, o_3\}, \{o_3, o_4\}\}$
- $ACL(o_1) = \{u_1\}$ $ACL(o_2) = \{\mathbf{u}_3\}$ $ACL(o_3) = \{\mathbf{u}_2\}$

 $ACL(o_4) = \{\mathbf{u}_3\}$

• policyLevel(read, o_1) = 2

policyLevel(write, o_1) = 0

policyLevel(read, o_2)= 2

 $policyLevel(write,o_2) = 1$

 $policyLevel(read,o_3) = 0$

 $policyLevel(write,o_3) = 0$

 $policyLevel(read,o_4) = 2$

 $policyLevel(write,o_4) = 1$

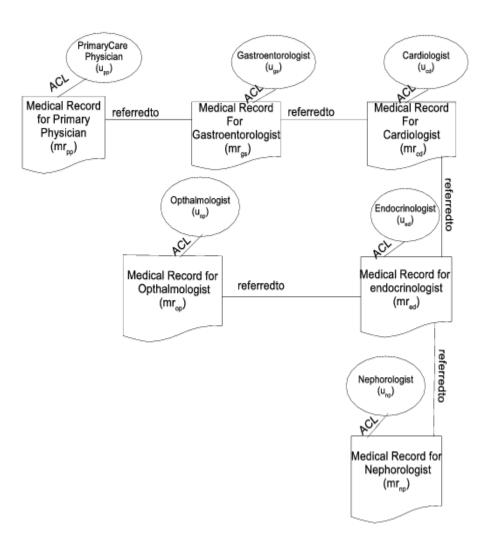
Sequence of operations and its outcome:

- read (u_1,o_3) , write (u_1,o_3) are denied
- read(u₂, o₁) is allowed, write(u₂, o₁) is denied
- read (u_1,o_4) , write (u_1,o_4) are denied



OOReBAC: Application





An OOReBAC Instantiation

- U = { upp, ugs, ucd, uop, ucd, unp }
 O = { mrpp, mrgs, mrcd, mrop, mred, mrnp }
 R = { {mrpp, mrgs}, {mrgs, mrcd}, {mrcd, mrcd, mrcd}, {mrcd, mrcd}, {mrop mred}, {mrnp, mred}}
 ACL(mrpp) = {upp},
 ACL(mrgs) = {ugs},
 ACL(mrcd) = {ucd},
 ACL(mrcd) = {ucd},
 ACL(mrcd) = {ucd},
 ACL(mrpp) = {unp}
 ACL(mrcd) = {ucd},
 ACL(mr
- policyLevel(read,mr $_{pp}$)= ∞ , policyLevel(write,mr $_{pp}$)=0, policyLevel(read,mr $_{gs}$)= ∞ , policyLevel(write,mr $_{gs}$)=0, policyLevel(read,mr $_{cd}$)= ∞ , policyLevel(write,mr $_{cd}$)=0, policyLevel(read,mr $_{op}$)= ∞ , policyLevel(write,mr $_{op}$)=0, policyLevel(read,mr $_{sd}$)= ∞ , policyLevel(write,mr $_{sd}$)=0, policyLevel(read,mr $_{np}$)= ∞ , policyLevel(write,mr $_{np}$)=0

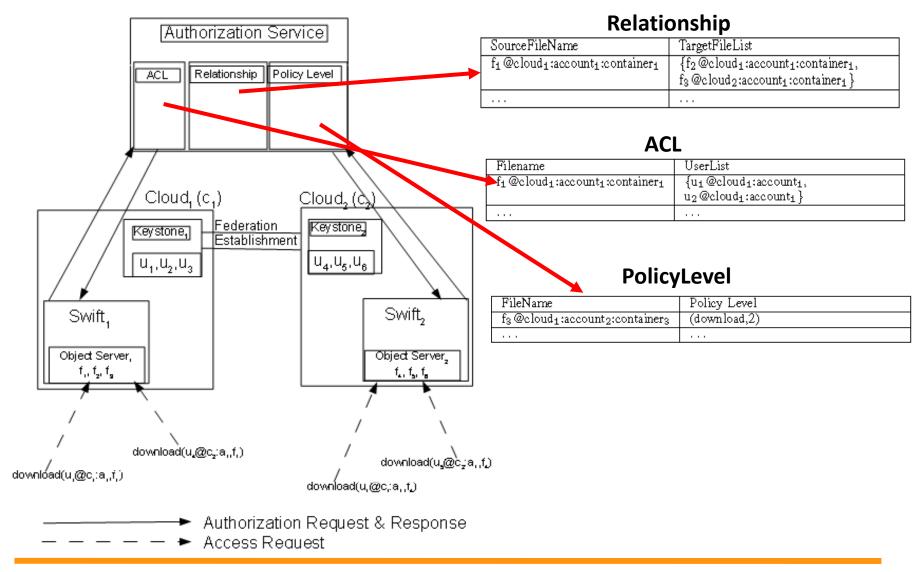
Sequence of Operations and Outcomes

- read(u_{np}, mr_{pp}): authorized
 read(u_{cd}, mr_{np}): authorized
- 3) write $(\mathbf{u}_{rxp}, \mathbf{mr}_{rxp})$: authorized
- 4) write $(\mathbf{u}_{rp}, \mathbf{mr}_{pp})$: denied
- 5) write $(\mathbf{u}_{rxp}, \mathbf{mr}_{pp})$: denied



Implementation: Openstack Object Storage (Swift)







Implementation



Functional Specification:

Functions	Conditions	Updates
	Administrative Actions	
CreateRelationship	$admin \in role(u) \land$	$RelationshipSet(filename_1) \hookrightarrow \{filename_2\}$
(u,filename ₁ ,filename ₂)	cloud(filename1) = cloud(u)	$RelationshipSet(filename_2) \cup= \{filename_1\}$
	∧ filename ₁ ∉ RelationshipSet(filename ₂)	
	∧ filename2 ∉ RelationshipSet(filename1)	
DeleteRelationship	$admin \in role(u) \land$	RelationshipSet(filename ₁) $\=$ {filename ₂ }
(u,filename ₁ ,filename ₂)	cloud(filename1) = cloud(u)	RelationshipSet(filename ₂) $=$ {filename ₁ }
	$filename_1 \in RelationshipSet(filename_2)$	
	\land filename ₂ \in RelationshipSet(filename ₁)	
IncludeAUserinACL	Role(u) ∈ Admin∧	ACLSet(filename1) ∪= {username1}
(u,filename ₁ ,username ₁)	$cloud(filename_1) = cloud(u) \land$	
	username₁ ∉ ACLSet(filename₁)	
ExcludeAUser FromACL	Role(u) ∈ Admin∧	$ACLSet(filename_1) = \{username_1\}$
(u,filename ₁ ,username ₁)	$cloud(filename_1) = cloud(u) \land$	
	$username_1 \in ACLSet(filename_2)$	
Configure Policy Level	Role(u)∈ Admin ∧	PolicyLevel(filename)= num
(u,filename,num)	$cloud(filename_1) = cloud(u)$	
	$ $ num $\leq \bigcirc $	
	Operational Command	
download	u∈ U	allow user u to download file filenamei
(u,filename ₁)	∧ authorize(u,filename ₁ ,G)	

Algorithm for Authorization

```
Algorithm 1 authorize(u,f,G)

if u in ACL(f) then
return true
else
policyLevel = policyLevel(f)
for depth limited search upto min(policyLevel, |O| - 1)
do
if if any of the file's ACL contains u then
return true
end if
end for
return false
end if
```

Conclusion and Future Work

- OOReBAC is the first attempt towards using object relationship independent of user in authorization policy specification and can only do where single type symmetric relationship is used.
- Limitations of OOReBAC:
 - Version Control system uses asymmetric relationship.
 - Object oriented Programming needs multiple Type asymmetric relationships.

We need to extend this model to accommodate multiple type asymmetric relationships to configure version control and object oriented system.





Questions?





